15.3 Implementing a Monitor System

Description

Here we will show how to implement a Monitor system using preemptive coroutines and Semaphore.

Below we show our first version of class Monitor:

```
class Monitor:
   entry:
     mutex.wait
     inner(entry)
     mutex.signal
%private
mutex: obj Semaphore(1)
```

A Monitor-object has a Semaphore attribute and a local method pattern, entry, which is to be used as a supermethod for methods in subclasses of Monitor. As can be seen, entry has an inner(entry); mutex.wait is executed before inner(entry) and mutex.signal is executed after inner(entry).

The Account class may be described as a subclass of Monitor in the following way:

```
class Account: Monitor
  deposit(amount: var float): entry
    balance := balance + amount
  withdraw(amount: var float): entry
    balance := balance - amount
  %private
  balance: var float
```

```
descriptor of
                                                    class Account: Monitor
class Monitor:
                                 Monitor
                                                        deposit(amount: var float
   entry:
                                                           balance := balance + a
      mutex.wait
                                                        withdraw(amount: var floa
      inner(entry)
                                                           balance := balance - a
      mutex.signal
                                    descriptor
                                                        %private
   %private
                                    of entry
                                                        balance: var float
   mutex: obj Semaphore(1)
```

If anAccount refers to an Account-object, then execution of anAccount.deposit(200) has the effect that the supermethod entry of deposit works like a wrapper around the statement balance := balance + amount. This ensures that mutex.signal is executed before the statement and mutex.signal is executed afterwards. The same is the case for execution of a anAccount.withDraw(300). All in all, using entry as a supermethod ensures that at most one deposit or withdraw may be executed at the same time, which guarantees exclusive acces to data-items within the Monitor object.

In section , Monitor is used as part of class that also defines a MonitorProcess class. We will show how to define such a system. It will contain the elements shown below:

```
class MonitorSystem
  class Monitor: ...
  class MonitorProcess: ...
```

```
%private
scheduler: obj ...
SQS: obj ProcessQueue
```

Class Monitor is defined as shown above. Class MonitorProcess is supposed to be a superclass of all parallel objects in the MonitorSystem. The scheduler object handles scheduling of MonitorProcess-objects and SQS is a queue of MonitorProcess-objects that are ready for being executed.

Class MonitorProcess may be defined as follows:

```
class MonitorProcess:
    start:
        status := ACTIVE
        SQS.insert(this(MonitorProcess))
    status: var integer
    inner(MonitorProcess)
    status := TERMINATED
```

A MonitorProcess has four attributes:

- A start method that sets the status of the MonitorProcess object to ACTIVE.
- An integer variable status holding the status of the MonitorProcess.
- An inner-statement that implies execution of items in a subclass of MonitorProcess.
- Finally a statement setting status to TERMINATED whereafter execution of the MonitorProcess object ends.

Next we describe the Scheduler object:

```
scheduler: obj
active: ref MonitorProcess
cycle
  active := SQS.next
  if (active <> none) :then
    active.attach(100)
    if (active.status = ACTIVE) then
        SQS.insert(active)
```

The scheduler has a data-item active that is a reference to the MonitorProcess object currently being executed.

Then it has a cycle-statement that forever executes:

- A reference to the next MonitorProcess in the queue SQS of active MonitorProcess-objects is assigned to active.
- If there are no references in SQS, active will get the datum none.
- If active is not none, the method active.attach(100) is invoked implying that execution of active is resumed. If active has not bee executed before, execution starts from the beginning of active; if active has been executed before ans execution has been suspended, execution is resumed after the point of suspension.
- Active will be preemptively suspended after 100 time units, but it may also terminate before the 100 time units have appeared.
- If active is preemptively suspended, active.status = ACTIVE and active is re-inserted into SQS.
- Otherwise active has terminated (status = TERMINATED) execution and thus not re-inserted into SQS.